

# Expressibility in the Lambda Calculus with letrec

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Utrecht University

TeReSe-Meeting

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# Which infinite $\lambda$ -terms are expressible finitely in $\lambda_{\text{letrec}}$ ?

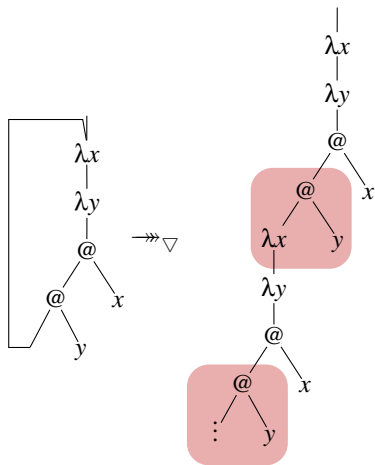
## Example

$\text{letrec } f = \lambda xy. f y x \text{ in } f \quad \rightsquigarrow_{\nabla} \quad \lambda xy. \lambda xy. (\lambda xy. (\dots y x) y x) y x$

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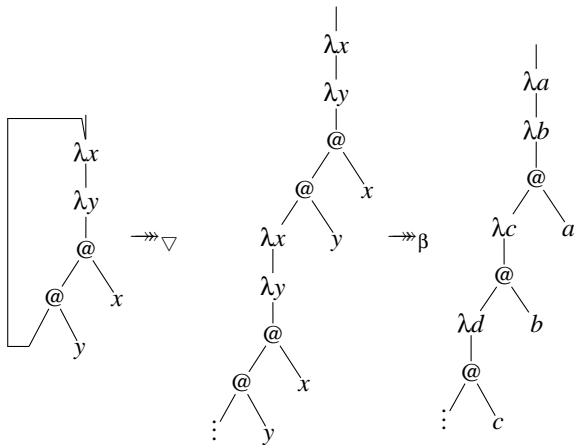
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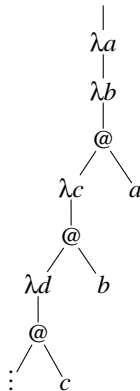




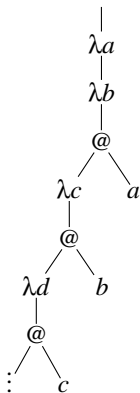


# What makes a $\lambda$ -terms inexpressible in $\lambda_{\text{letrec}}$ ?

## Example



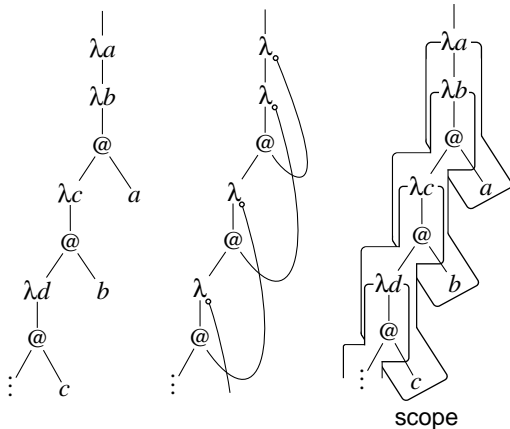
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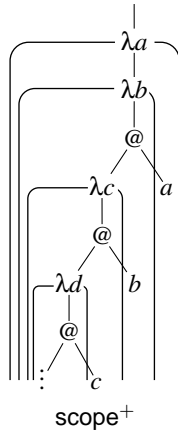
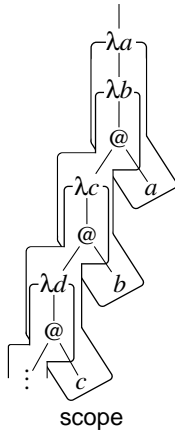
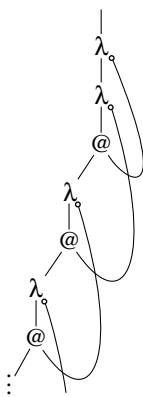
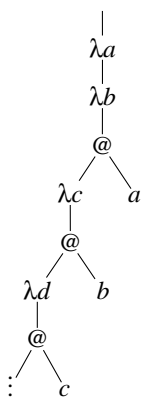




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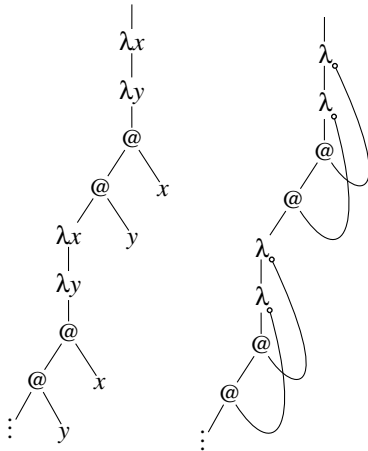


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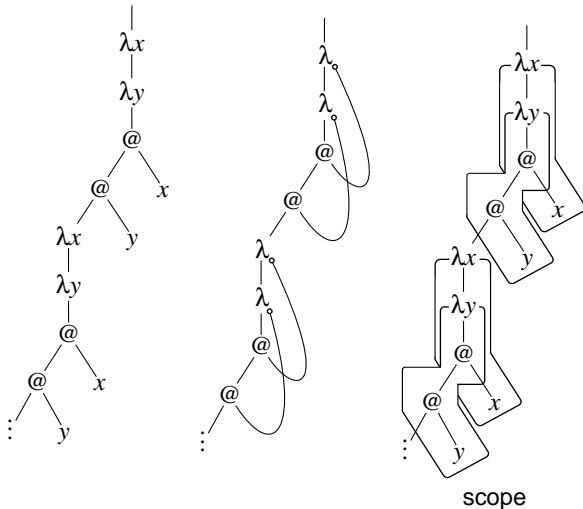




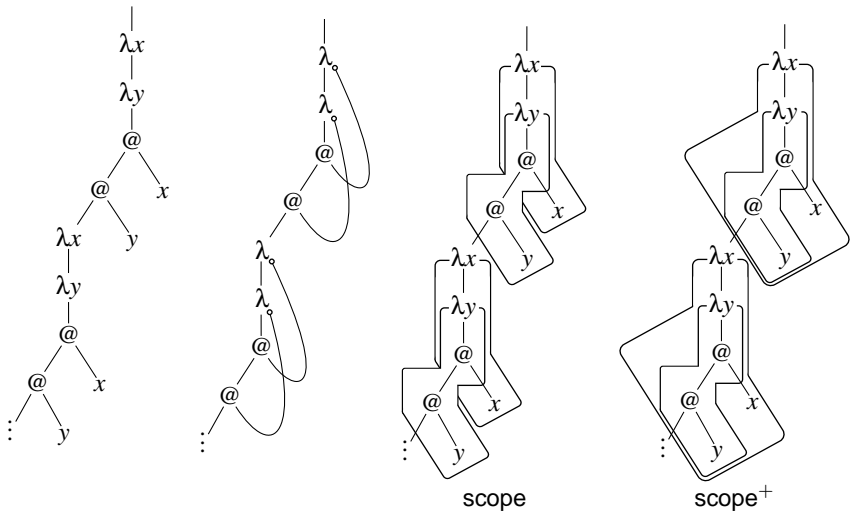
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# A CRS for observing infinite $\lambda$ -terms: **Reg<sup>+</sup>**

$$(\rho^{\circledast i}) : \quad (\lambda x_1 \dots x_n) M_0 M_1 \rightarrow (\lambda x_1 \dots x_n) M_i \quad (i \in \{0, 1\})$$

$$(\rho^\lambda) : \quad (\lambda x_1 \dots x_n) \lambda x_{n+1}. M_0 \rightarrow (\lambda x_1 \dots x_{n+1}) M_0$$

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$$(\rho^{\odot i}) : \text{pre}_n([x_1 \dots x_n] \text{app}(Z_0(\vec{x}), Z_1(\vec{x}))) \rightarrow \text{pre}_n([x_1 \dots x_n] Z_i(\vec{x})) \quad (i \in \{0, 1\})$$

$$(\rho^\lambda) : \text{pre}_n([x_1 \dots x_n] \text{abs}([x_{n+1}] Z(\vec{x}))) \rightarrow \text{pre}_n([x_1 \dots x_{n+1}] Z(\vec{x}))$$

$$(\rho^S) : \text{pre}_{n+1}([x_1 \dots x_{n+1}] Z(x_1, \dots, x_n)) \rightarrow \text{pre}_n([x_1 \dots x_n] Z(x_1, \dots, x_n))$$

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Example  $(\lambda x. \lambda y. x x y)$

$() \lambda x. \lambda y. x x y$

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$$\begin{aligned} & () \lambda x. \lambda y. x x y \rightarrow_\lambda \\ & (\lambda x) \lambda y. x x y \end{aligned}$$

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$$\begin{aligned}
 () \lambda x. \lambda y. x x y &\rightarrow_\lambda \\
 (\lambda x) \lambda y. x x y &\rightarrow_\lambda \\
 (\lambda x y) x x y &\rightarrow_{\textcircled{1}} \\
 (\lambda x y) y &
 \end{aligned}$$

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 (\lambda x) \lambda y. x x y \rightarrow_\lambda & (\lambda x y) x x y \rightarrow_{\textcircled{0}} \\
 (\lambda x y) x x y \rightarrow_{\textcircled{1}} & (\lambda x y) x x \rightarrow_{\textcircled{1}} \\
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 (\lambda x) \lambda y. x x y \rightarrow_\lambda & (\lambda x y) x x y \rightarrow_{@_0} \\
 (\lambda x y) x x y \rightarrow_{@_1} & (\lambda x y) x x \rightarrow_{@_1} \\
 (\lambda x y) y & (\lambda x y) x \rightarrow_S \\
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 (\lambda x y) x x y \rightarrow_{\circledast_1} & (\lambda x y) x x \rightarrow_{\circledast_1} & \\
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$(\lambda x) \lambda y. x x y \rightarrow_\lambda$	$(\lambda x y) x x y \rightarrow_{@_0}$	$(\lambda x y) x x y$
$(\lambda x y) x x y \rightarrow_{@_1}$	$(\lambda x y) x x \rightarrow_{@_1}$	
$(\lambda x y) y$	$(\lambda x y) x \rightarrow_S$	
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$(\lambda x y) x x y \rightarrow_{@_1}$	$(\lambda x y) x x \rightarrow_{@_1}$	$(\lambda x y) x x$
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$(\lambda x y) x x y \rightarrow_{@_1}$	$(\lambda x y) x x \rightarrow_{@_1}$	$(\lambda x y) x x \rightarrow_S$
$(\lambda x y) y$	$(\lambda x y) x \rightarrow_S$	$(\lambda x) x x$
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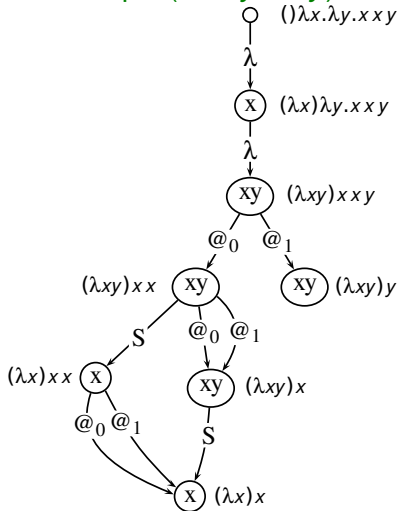
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$(\lambda x y) y$	$(\lambda x y) x \rightarrow_S$	$(\lambda x) x x \rightarrow_{@_1}$
	$(\lambda x) x$	$(\lambda x) x$

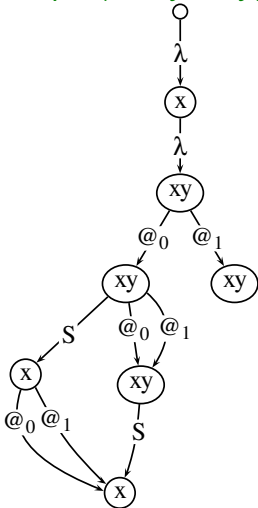
# scope<sup>+</sup>-delimiting strategies for *Reg*<sup>+</sup>

Example  $(\lambda x.\lambda y.x x y)$



# scope<sup>+</sup>-delimiting strategies for *Reg*<sup>+</sup>

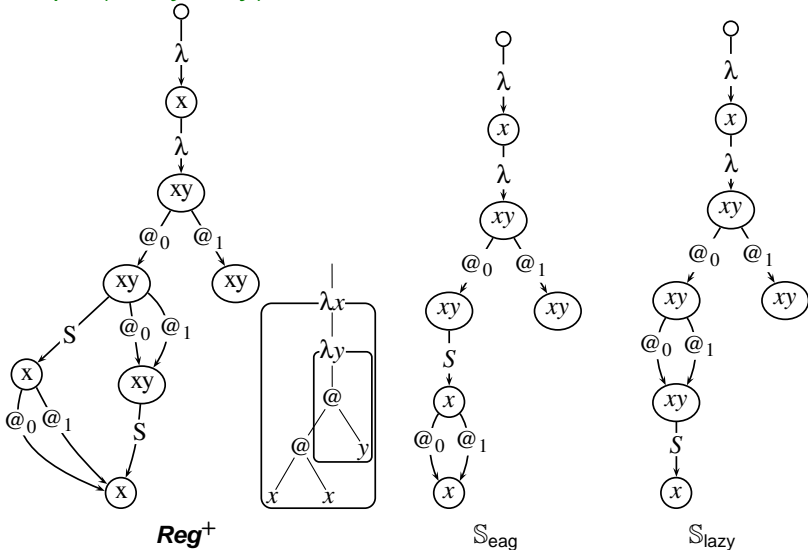
Example  $(\lambda x.\lambda y.x x y)$





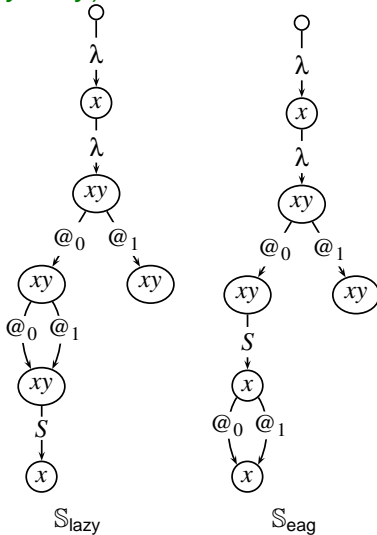
# scope<sup>+</sup>-delimiting strategies for *Reg*<sup>+</sup>

Example  $(\lambda x.\lambda y.x x y)$



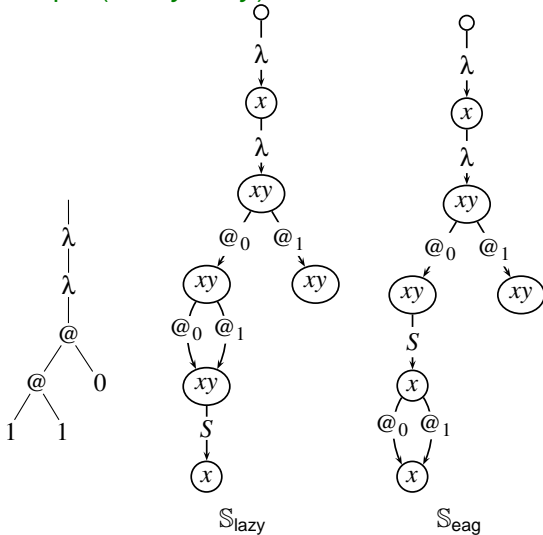
# Connection to the de-Bruijn notation

Example  $(\lambda x. \lambda y. xxy)$



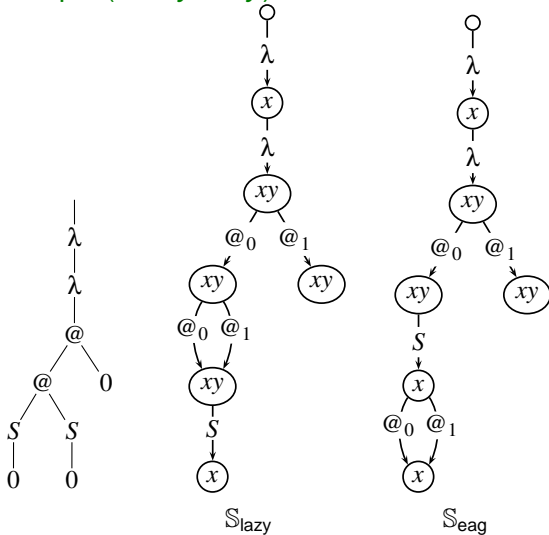
# Connection to the de-Bruijn notation

Example  $(\lambda x. \lambda y. xxy)$



# Connection to the de-Brujin notation

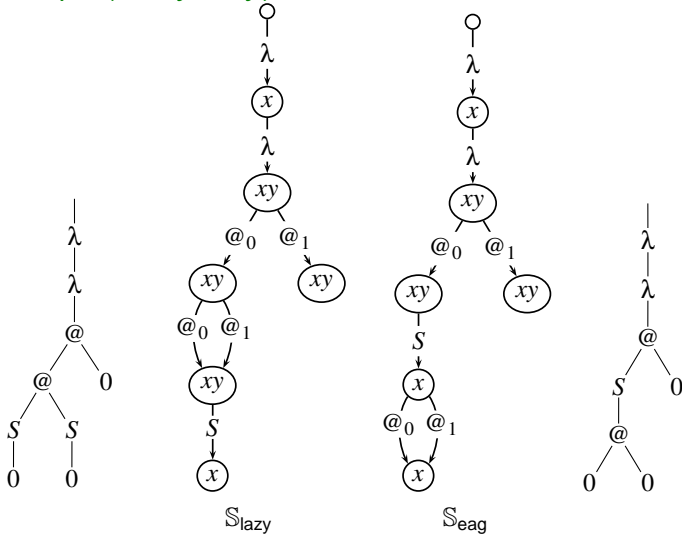
Example  $(\lambda x. \lambda y. x x y)$





# Connection to the de-Bruijn notation

Example  $(\lambda x. \lambda y. xxy)$



# Strong regularity

## Definition ( $\mathbb{S}$ -regularity)

An infinite  $\lambda$ -term  $M$  is  *$\mathbb{S}$ -regular* the set of  $\rightarrow_{\mathbb{S}}$ -reducts of  $M$  is finite.

# Strong regularity

## Definition ( $\mathbb{S}$ -regularity)

An infinite  $\lambda$ -term  $M$  is  *$\mathbb{S}$ -regular* if the set of  $\rightarrow_{\mathbb{S}}$ -reducts of  $M$  is finite.

## Definition (strongly regular infinite $\lambda$ -terms)

An infinite  $\lambda$ -term  $M$  is *strongly regular* if there exists a scope<sup>+</sup>-delimiting strategy  $\mathbb{S}^+$  for  $\text{Reg}^+$  such that  $M$  is  $\mathbb{S}^+$ -regular.

# $\lambda_{\text{letrec}}$ -expressibility

## Definition

A  $\lambda_{\text{letrec}}$ -term  $L$  *expresses* an infinite  $\lambda$ -term  $M$  if  $L \rightarrow_{\nabla}^{\omega} M$ .

## Definition

$M$  is  *$\lambda_{\text{letrec}}$ -expressible* if there is a  $\lambda_{\text{letrec}}$ -term which expresses  $M$ .

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## Definition

$M$  is  $\lambda_{\text{letrec}}$ -*expressible* if there is a  $\lambda_{\text{letrec}}$ -term which expresses  $M$ .

$$(\rho_{\nabla}^{\circ}) : \quad \text{letrec } B \text{ in } L_0 L_1 \rightarrow (\text{letrec } B \text{ in } L_0) (\text{letrec } B \text{ in } L_1)$$

$$(\rho_{\nabla}^{\lambda}) : \quad \text{letrec } B \text{ in } \lambda x. L_0 \rightarrow \lambda x. \text{letrec } B \text{ in } L_0$$

$$(\rho_{\nabla}^{\text{letrec}}) : \quad \text{letrec } B_0 \text{ in letrec } B_1 \text{ in } L \rightarrow \text{letrec } B_0, B_1 \text{ in } L$$

$$(\rho_{\nabla}^{\text{rec}}) : \quad \text{letrec } B \text{ in } f_i \rightarrow \text{letrec } B \text{ in } L_i$$

(if  $B$  is  $f_1 = L_1 \dots f_n = L_n$ )

$$(\rho_{\nabla}^{\text{red}}) : \quad \text{letrec } f_1 = L_1 \dots f_n = L_n \text{ in } L \rightarrow \text{letrec } f_{j_1} = L_{j_1} \dots f_{j_{n'}} = L_{j_{n'}} \text{ in } L$$

(if  $f_{j_1}, \dots, f_{j_{n'}}$  are the recursion variables reachable from  $L$ )

$$(\rho_{\nabla}^{\text{nil}}) : \quad \text{letrec in } L \rightarrow L$$

# Strong regularity and $\lambda_{\text{letrec}}$ -expressibility

## Theorem

*An infinite  $\lambda$ -term is  $\lambda_{\text{letrec}}$ -expressible if and only if it is strongly regular.*

# Strong regularity

Example ( $M = \lambda xy. M y x$ )

$$\frac{S_{\text{lazy}}^+}{()M}$$

# Strong regularity

Example ( $M = \lambda xy. M y x$ )

$$\frac{S_{\text{lazy}}^+}{\begin{array}{l} ()M \\ ()\lambda xy. M y x \end{array}} =$$



# Strong regularity

Example ( $M = \lambda xy.M y x$ )

$$\begin{array}{l}
 \mathbb{S}_{\text{lazy}}^+ \\
 \hline
 ()M \quad \quad \quad = \\
 ()\lambda xy.M y x \quad \quad \rightarrow_{\lambda} \\
 (\lambda x)\lambda y.M y x
 \end{array}$$

# Strong regularity

Example ( $M = \lambda xy. My x$ )

$S_{\text{lazy}}^+$	
()	=
$M$	
$(\lambda xy. My x)$	$\rightarrow_\lambda$
$(\lambda x)\lambda y. My x$	$\rightarrow_\lambda$
$(\lambda xy) My x$	

# Strong regularity

Example ( $M = \lambda xy. My x$ )

$S_{\text{lazy}}^+$		
$()M$		$=$
$()\lambda xy. My x$		$\rightarrow_{\lambda}$
$(\lambda x)\lambda y. My x$		$\rightarrow_{\lambda}$
$(\lambda xy)My x$		$\rightarrow_{@_0}$
$(\lambda xy)My$		

# Strong regularity

Example ( $M = \lambda xy.M y x$ )

$S_{\text{lazy}}^+$	
$()M$	$=$
$()\lambda xy.M y x$	$\rightarrow_{\lambda}$
$(\lambda x)\lambda y.M y x$	$\rightarrow_{\lambda}$
$(\lambda xy)M y x$	$\rightarrow_{@_0}$
$(\lambda xy)M y$	$\rightarrow_{@_0}$
$(\lambda xy)M$	

# Strong regularity

Example ( $M = \lambda xy. My x$ )

$S_{\text{lazy}}^+$		
	$()M$	$=$
	$()\lambda xy. My x$	$\rightarrow_{\lambda}$
	$(\lambda x)\lambda y. My x$	$\rightarrow_{\lambda}$
	$(\lambda xy)My x$	$\rightarrow_{@_0}$
	$(\lambda xy)My$	$\rightarrow_{@_0}$
	$(\lambda xy)M$	$=$
	$(\lambda xy)\lambda x' y'. My' x'$	

# Strong regularity

Example ( $M = \lambda xy. My x$ )

$S_{\text{lazy}}^+$		
$()M$		=
$()\lambda xy. My x$		$\rightarrow_\lambda$
$(\lambda x)\lambda y. My x$		$\rightarrow_\lambda$
$(\lambda xy)My x$		$\rightarrow_{@_0}$
$(\lambda xy)My$		$\rightarrow_{@_0}$
$(\lambda xy)M$		=
$(\lambda xy)\lambda x' y'. My' x'$		$\rightarrow_\lambda$
$(\lambda xyx')\lambda y'. My' x'$		

# Strong regularity

Example ( $M = \lambda xy. My x$ )

$S_{\text{lazy}}^+$	
$()M$	=
$()\lambda xy. My x$	$\rightarrow_\lambda$
$(\lambda x)\lambda y. My x$	$\rightarrow_\lambda$
$(\lambda xy)My x$	$\rightarrow_{@_0}$
$(\lambda xy)My$	$\rightarrow_{@_0}$
$(\lambda xy)M$	=
$(\lambda xy)\lambda x' y'. My' x'$	$\rightarrow_\lambda$
$(\lambda xyx')\lambda y'. My' x'$	$\rightarrow_\lambda$
$(\lambda xyx' y')My' x'$	

# Strong regularity

Example ( $M = \lambda xy. My x$ )

$S_{\text{lazy}}^+$	
$()M$	=
$()\lambda xy. My x$	$\rightarrow_{\lambda}$
$(\lambda x)\lambda y. My x$	$\rightarrow_{\lambda}$
$(\lambda xy)My x$	$\rightarrow_{@_0}$
$(\lambda xy)My$	$\rightarrow_{@_0}$
$(\lambda xy)M$	=
$(\lambda xy)\lambda x' y'. My' x'$	$\rightarrow_{\lambda}$
$(\lambda xyx')\lambda y'. My' x'$	$\rightarrow_{\lambda}$
$(\lambda xyx' y')My' x'$	$\rightarrow_{@_0}$
$(\lambda xyx' y')My'$	



# Strong regularity

Example ( $M = \lambda xy. My x$ )

$S_{\text{lazy}}^+$	
$()M$	$=$
$()\lambda xy. My x$	$\rightarrow_{\lambda}$
$(\lambda x)\lambda y. My x$	$\rightarrow_{\lambda}$
$(\lambda xy)My x$	$\rightarrow_{@_0}$
$(\lambda xy)My$	$\rightarrow_{@_0}$
$(\lambda xy)M$	$=$
$(\lambda xy)\lambda x' y'. My' x'$	$\rightarrow_{\lambda}$
$(\lambda xyx')\lambda y'. My' x'$	$\rightarrow_{\lambda}$
$(\lambda xyx' y')My' x'$	$\rightarrow_{@_0}$
$(\lambda xyx' y')My'$	$\rightarrow_{@_0}$
$(\lambda xyx' y')M$	$\dots$

# Strong regularity

Example ( $M = \lambda xy. Myx$ )

$S_{\text{lazy}}^+$			
()	$M$	=	$S_{\text{eag}}^+$
	$(\lambda xy. Myx)$	$\rightarrow_{\lambda}$	()
	$(\lambda x)\lambda y. Myx$	$\rightarrow_{\lambda}$	$M$
	$(\lambda xy)Myx$	$\rightarrow_{@_0}$	
	$(\lambda xy)My$	$\rightarrow_{@_0}$	
	$(\lambda xy)M$	=	
	$(\lambda xy)\lambda x'y'. My'x'$	$\rightarrow_{\lambda}$	
	$(\lambda xyx')\lambda y'. My'x'$	$\rightarrow_{\lambda}$	
	$(\lambda xyx'y')My'x'$	$\rightarrow_{@_0}$	
	$(\lambda xyx'y')My'$	$\rightarrow_{@_0}$	
	$(\lambda xyx'y')M$	...	

# Strong regularity

Example ( $M = \lambda xy. Myx$ )

$$\begin{array}{l}
 \frac{}{S_{\text{lazy}}^+} \\
 \hline
 ()M \quad = \\
 ()\lambda xy. Myx \quad \rightarrow_{\lambda} \\
 (\lambda x)\lambda y. Myx \quad \rightarrow_{\lambda} \\
 (\lambda xy)Myx \quad \rightarrow_{@_0} \\
 (\lambda xy)My \quad \rightarrow_{@_0} \\
 (\lambda xy)M \quad = \\
 (\lambda xy)\lambda x'y'. My'x' \quad \rightarrow_{\lambda} \\
 (\lambda xyx')\lambda y'. My'x' \quad \rightarrow_{\lambda} \\
 (\lambda xyx'y')My'x' \quad \rightarrow_{@_0} \\
 (\lambda xyx'y')My' \quad \rightarrow_{@_0} \\
 (\lambda xyx'y')M \quad \dots
 \end{array}
 \quad = \quad
 \frac{}{S_{\text{eag}}^+}
 \begin{array}{l}
 ()M \quad = \\
 ()\lambda xy. Myx
 \end{array}$$

# Strong regularity

Example ( $M = \lambda xy.Myx$ )

$S_{\text{lazy}}^+$			$S_{\text{eag}}^+$	
$()M$	$=$		$()M$	$=$
$()\lambda xy.Myx$	$\rightarrow\lambda$		$()\lambda xy.Myx$	$\rightarrow\lambda$
$(\lambda x)\lambda y.Myx$	$\rightarrow\lambda$		$(\lambda x)\lambda y.Myx$	
$(\lambda xy)Myx$	$\rightarrow@_0$			
$(\lambda xy)My$	$\rightarrow@_0$			
$(\lambda xy)M$	$=$			
$(\lambda xy)\lambda x'y'.My'x'$	$\rightarrow\lambda$			
$(\lambda xyx')\lambda y'.My'x'$	$\rightarrow\lambda$			
$(\lambda xyx'y')My'x'$	$\rightarrow@_0$			
$(\lambda xyx'y')My'$	$\rightarrow@_0$			
$(\lambda xyx'y')M$	$\dots$			

# Strong regularity

Example ( $M = \lambda xy.Myx$ )

$\frac{S_{\text{lazy}}^+}{\text{---}}$			$\frac{S_{\text{eag}}^+}{\text{---}}$		
$()M$	=		$()M$	=	
$()\lambda xy.Myx$	$\rightarrow_\lambda$		$()\lambda xy.Myx$	$\rightarrow_\lambda$	
$(\lambda x)\lambda y.Myx$	$\rightarrow_\lambda$		$(\lambda x)\lambda y.Myx$	$\rightarrow_\lambda$	
$(\lambda xy)Myx$	$\rightarrow_{@_0}$		$(\lambda xy)Myx$		
$(\lambda xy)My$	$\rightarrow_{@_0}$				
$(\lambda xy)M$	=				
$(\lambda xy)\lambda x'y'.My'x'$	$\rightarrow_\lambda$				
$(\lambda xyx')\lambda y'.My'x'$	$\rightarrow_\lambda$				
$(\lambda xyx'y')My'x'$	$\rightarrow_{@_0}$				
$(\lambda xyx'y')My'$	$\rightarrow_{@_0}$				
$(\lambda xyx'y')M$	...				

# Strong regularity

Example ( $M = \lambda xy. Myx$ )

$S_{\text{lazy}}^+$			$S_{\text{eag}}^+$	
$()M$	=		$()M$	=
$()\lambda xy. Myx$	$\rightarrow_\lambda$		$()\lambda xy. Myx$	$\rightarrow_\lambda$
$(\lambda x)\lambda y. Myx$	$\rightarrow_\lambda$		$(\lambda x)\lambda y. Myx$	$\rightarrow_\lambda$
$(\lambda xy)Myx$	$\rightarrow_{@_0}$		$(\lambda xy)Myx$	$\rightarrow_{@_0}$
$(\lambda xy)My$	$\rightarrow_{@_0}$		$(\lambda xy)My$	
$(\lambda xy)M$	=		$(\lambda xy)My$	
$(\lambda xy)\lambda x'y'. My'x'$	$\rightarrow_\lambda$			
$(\lambda xyx')\lambda y'. My'x'$	$\rightarrow_\lambda$			
$(\lambda xyx'y')My'x'$	$\rightarrow_{@_0}$			
$(\lambda xyx'y')My'$	$\rightarrow_{@_0}$			
$(\lambda xyx'y')M$	...			

# Strong regularity

Example ( $M = \lambda xy.Myx$ )

$S_{\text{lazy}}^+$			$S_{\text{eag}}^+$	
$()M$	=		$()M$	=
$()\lambda xy.Myx$	$\rightarrow\lambda$		$()\lambda xy.Myx$	$\rightarrow\lambda$
$(\lambda x)\lambda y.Myx$	$\rightarrow\lambda$		$(\lambda x)\lambda y.Myx$	$\rightarrow\lambda$
$(\lambda xy)Myx$	$\rightarrow@_0$		$(\lambda xy)Myx$	$\rightarrow@_0$
$(\lambda xy)My$	$\rightarrow@_0$		$(\lambda xy)My$	$\rightarrow@_0$
$(\lambda xy)M$	=		$(\lambda xy)M$	$\rightarrow@_0$
$(\lambda xy)\lambda x'y'.My'x'$	$\rightarrow\lambda$		$(\lambda xy)M$	
$(\lambda xyx')\lambda y'.My'x'$	$\rightarrow\lambda$			
$(\lambda xyx'y')My'x'$	$\rightarrow@_0$			
$(\lambda xyx'y')My'$	$\rightarrow@_0$			
$(\lambda xyx'y')M$	...			

# Strong regularity

Example ( $M = \lambda xy.Myx$ )

$S_{\text{lazy}}^+$			$S_{\text{eag}}^+$	
$()M$	=		$()M$	=
$()\lambda xy.Myx$	$\rightarrow\lambda$		$()\lambda xy.Myx$	$\rightarrow\lambda$
$(\lambda x)\lambda y.Myx$	$\rightarrow\lambda$		$(\lambda x)\lambda y.Myx$	$\rightarrow\lambda$
$(\lambda xy)Myx$	$\rightarrow@_0$		$(\lambda xy)Myx$	$\rightarrow@_0$
$(\lambda xy)My$	$\rightarrow@_0$		$(\lambda xy)My$	$\rightarrow@_0$
$(\lambda xy)M$	=		$(\lambda xy)M$	$\rightarrow S$
$(\lambda xy)\lambda x'y'.My'x'$	$\rightarrow\lambda$		$(\lambda x)M$	
$(\lambda xyx')\lambda y'.My'x'$	$\rightarrow\lambda$			
$(\lambda xyx'y')My'x'$	$\rightarrow@_0$			
$(\lambda xyx'y')My'$	$\rightarrow@_0$			
$(\lambda xyx'y')M$	...			



# Strong regularity

Example ( $M = \lambda xy.Myx$ )

$S_{\text{lazy}}^+$			$S_{\text{eag}}^+$	
$()M$	=		$()M$	=
$()\lambda xy.Myx$	$\rightarrow\lambda$		$()\lambda xy.Myx$	$\rightarrow\lambda$
$(\lambda x)\lambda y.Myx$	$\rightarrow\lambda$		$(\lambda x)\lambda y.Myx$	$\rightarrow\lambda$
$(\lambda xy)Myx$	$\rightarrow@_0$		$(\lambda xy)Myx$	$\rightarrow@_0$
$(\lambda xy)My$	$\rightarrow@_0$		$(\lambda xy)My$	$\rightarrow@_0$
$(\lambda xy)M$	=		$(\lambda xy)M$	$\rightarrow S$
$(\lambda xy)\lambda x'y'.My'x'$	$\rightarrow\lambda$		$(\lambda x)M$	$\rightarrow S$
$(\lambda xyx')\lambda y'.My'x'$	$\rightarrow\lambda$		$()M$	...
$(\lambda xyx'y')My'x'$	$\rightarrow@_0$			
$(\lambda xyx'y')My'$	$\rightarrow@_0$			
$(\lambda xyx'y')M$	...			

# Strong regularity and $\mathbb{S}_{\text{eag}}^+$

## Proposition

$M$  is *strongly regular* if and only if  $M$  is  $\mathbb{S}_{\text{eag}}^+$ -regular.

## Proof.

$\Leftarrow$ : by definition.

$\Rightarrow$ : follows from  $\rightarrow_{\mathbb{S}_{\text{eag}}^+} \subseteq \rightarrow_{\mathbb{S}^+} \cdot \rightarrow_{\mathbb{S}}$

# Strong regularity and $\mathbb{S}_{\text{eag}}^+$

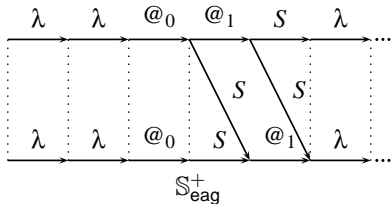
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$M$  is *strongly regular* if and only if  $M$  is  $\mathbb{S}_{\text{eag}}^+$ -regular.

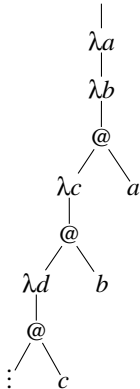
## Proof.

$\Leftarrow$ : by definition.

$\Rightarrow$ : follows from  $\rightarrow_{\mathbb{S}_{\text{eag}}^+} \subseteq \rightarrow_{\mathbb{S}^+} \cdot \rightarrow_{\mathbb{S}}$

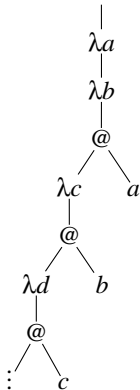


# Strong regularity and $S_{eag}^+$



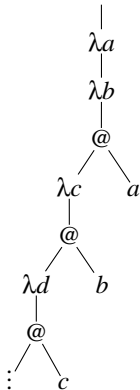
$$() \lambda a. \lambda b. \dots \rightarrow_{\lambda} (\lambda a) \lambda b. \lambda c. \dots a$$

# Strong regularity and $S_{eag}^+$



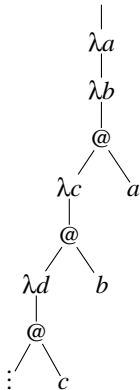
$$\begin{aligned}
 ()\lambda a.\lambda b.\dots &\rightarrow_{\lambda} (\lambda a)\lambda b.\lambda c.\dots a \\
 &\rightarrow_{\lambda} (\lambda ab)\lambda c.\dots a
 \end{aligned}$$

# Strong regularity and $S_{eag}^+$



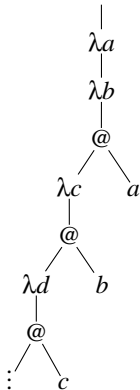
$$\begin{aligned}
 ()\lambda a.\lambda b.\dots &\rightarrow_{\lambda} (\lambda a)\lambda b.\lambda c.\dots a \\
 &\rightarrow_{\lambda} (\lambda ab)\lambda c.\dots a \\
 &\rightarrow_{@_0} (\lambda ab)\lambda c.\lambda d.\dots b
 \end{aligned}$$

# Strong regularity and $S_{eag}^+$



$$\begin{array}{ll}
 ()\lambda a.\lambda b.\dots & \rightarrow_{\lambda} (\lambda a)\lambda b.\lambda c.\dots a \\
 & \rightarrow_{\lambda} (\lambda ab)\lambda c.\dots a \\
 & \rightarrow_{@_0} (\lambda ab)\lambda c.\lambda d.\dots b \\
 & \rightarrow_{\lambda} (\lambda abc)\lambda d.\dots b
 \end{array}$$

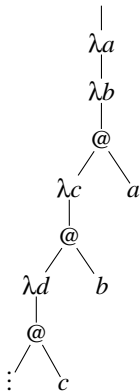
# Strong regularity and $S_{eag}^+$



$$\begin{array}{ll}
 ()\lambda a.\lambda b.\dots & \rightarrow_{\lambda} (\lambda a)\lambda b.\lambda c.\dots a \\
 & \rightarrow_{\lambda} (\lambda ab)\lambda c.\dots a \\
 & \rightarrow_{@_0} (\lambda ab)\lambda c.\lambda d.\dots b \\
 & \rightarrow_{\lambda} (\lambda abc)\lambda d.\dots b \\
 & \rightarrow_{@_0} (\lambda abc)\lambda d.\lambda e.\dots c
 \end{array}$$

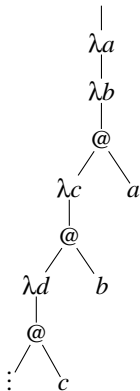


# Strong regularity and $S_{\text{eag}}^+$



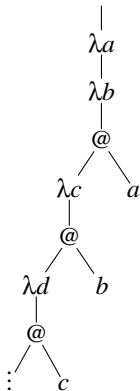
$$\begin{array}{ll}
 ()\lambda a.\lambda b.\dots & \rightarrow_{\lambda} (\lambda a)\lambda b.\lambda c.\dots a \\
 & \rightarrow_{\lambda} (\lambda ab)\lambda c.\dots a \\
 & \rightarrow_{@_0} (\lambda ab)\lambda c.\lambda d.\dots b \\
 & \rightarrow_{\lambda} (\lambda abc)\lambda d.\dots b \\
 & \rightarrow_{@_0} (\lambda abc)\lambda d.\lambda e.\dots c \\
 & \rightarrow_{\lambda} (\lambda abcd)\lambda e.\dots c
 \end{array}$$

# Strong regularity and $S_{\text{eag}}^+$



$$\begin{array}{ll}
 ()\lambda a.\lambda b.\dots & \rightarrow_{\lambda} (\lambda a)\lambda b.\lambda c.\dots a \\
 & \rightarrow_{\lambda} (\lambda ab)\lambda c.\dots a \\
 & \rightarrow_{@_0} (\lambda ab)\lambda c.\lambda d.\dots b \\
 & \rightarrow_{\lambda} (\lambda abc)\lambda d.\dots b \\
 & \rightarrow_{@_0} (\lambda abc)\lambda d.\lambda e.\dots c \\
 & \rightarrow_{\lambda} (\lambda abcd)\lambda e.\dots c \\
 & \rightarrow_{@_0} (\lambda abcd)\lambda e.\lambda f.\dots d
 \end{array}$$

# Strong regularity and $S_{eag}^+$



$$\begin{array}{ll}
 ()\lambda a.\lambda b.\dots & \rightarrow_{\lambda} (\lambda a)\lambda b.\lambda c.\dots a \\
 & \rightarrow_{\lambda} (\lambda ab)\lambda c.\dots a \\
 & \rightarrow_{@_0} (\lambda ab)\lambda c.\lambda d.\dots b \\
 & \rightarrow_{\lambda} (\lambda abc)\lambda d.\dots b \\
 & \rightarrow_{@_0} (\lambda abc)\lambda d.\lambda e.\dots c \\
 & \rightarrow_{\lambda} (\lambda abcd)\lambda e.\dots c \\
 & \rightarrow_{@_0} (\lambda abcd)\lambda e.\lambda f.\dots d \\
 & \rightarrow_{\lambda} (\lambda abcde)\lambda f.\dots d \\
 & \dots
 \end{array}$$

# Proof system $\mathbf{Reg}^+$ for strong regularity

$$\frac{}{(\lambda \vec{x} y) y} 0$$

$$\frac{(\lambda \vec{x} y) M_0}{(\lambda \vec{x}) \lambda y. M_0} \lambda$$

$$\frac{(\lambda \vec{x}) M_0 \quad (\lambda \vec{x}) M_1}{(\lambda \vec{x}) M_0 M_1} @$$

$$\frac{(\lambda x_1 \dots x_{n-1}) M}{(\lambda x_1 \dots x_n) M} S \quad \text{(if the binding } \lambda x_n \text{ is vacuous)}$$

$$[(\lambda \vec{x}) M]^u$$

$$\frac{\mathcal{D}_0 \quad (\lambda \vec{x}) M}{(\lambda \vec{x}) M} \text{FIX}, u \quad \text{(if } |\mathcal{D}_0| \geq 1)$$

# Proof system $\mathbf{Reg}^+$ for strong regularity

$$\begin{array}{c}
 \frac{}{(\lambda \vec{x} y) y} \text{0} \qquad \frac{(\lambda \vec{x} y) M_0}{(\lambda \vec{x}) \lambda y. M_0} \lambda \qquad \frac{(\lambda \vec{x}) M_0 \quad (\lambda \vec{x}) M_1}{(\lambda \vec{x}) M_0 M_1} \textcircled{c} \\
 \\
 \frac{(\lambda x_1 \dots x_{n-1}) M}{(\lambda x_1 \dots x_n) M} \text{S} \quad \text{(if the binding } \lambda x_n \text{ is vacuous)} \qquad \frac{[(\lambda \vec{x}) M]^u \quad \mathcal{D}_0}{(\lambda \vec{x}) M} \text{FIX, } u \text{ (if } |\mathcal{D}_0| \geq 1)
 \end{array}$$

$$(\rho^{\textcircled{i}}) : \quad (\lambda x_1 \dots x_n) M_0 M_1 \rightarrow (\lambda x_1 \dots x_n) M_i \quad (i \in \{0, 1\})$$

$$(\rho^\lambda) : \quad (\lambda x_1 \dots x_n) \lambda x_{n+1}. M_0 \rightarrow (\lambda x_1 \dots x_{n+1}) M_0$$

$$(\rho^{\text{S}}) : \quad (\lambda x_1 \dots x_{n+1}) M_0 \rightarrow (\lambda x_1 \dots x_n) M_0 \quad (\text{if } \lambda x_{n+1} \text{ is vacuous)}$$

# Proof system $\mathbf{Reg}^+$ for strong regularity

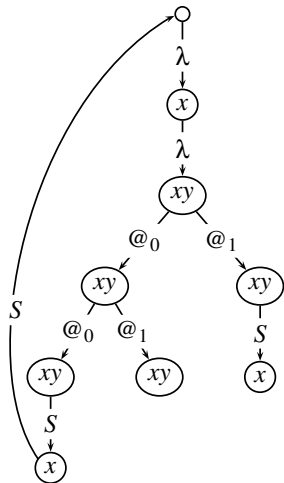
Example ( $M = \lambda xy. M x y$ )

$$\begin{array}{c}
 \frac{((\ ))M^u}{(\lambda x)M} \text{ S} \quad \frac{}{(\lambda x)x} \text{ 0} \\
 \hline
 \quad \quad \quad \text{©} \\
 \frac{(\lambda x)Mx}{(\lambda xy)Mx} \text{ S} \quad \frac{}{(\lambda xy)y} \text{ 0} \\
 \hline
 \quad \quad \quad \text{©} \\
 \frac{(\lambda xy)Mxy}{(\lambda x)\lambda y.Mxy} \lambda \\
 \hline
 \quad \quad \quad \lambda \\
 \frac{(\ )M}{(\ )M} \text{ FIX, } u
 \end{array}$$

# Proof system $\mathbf{Reg}^+$ for strong regularity

Example ( $M = \lambda xy. Mxy$ )

$$\begin{array}{c}
 \frac{((\ ))M^u}{(\lambda x)M} S \quad \frac{}{(\lambda x)x} 0}{\frac{}{(\lambda x)Mx} 0} @ \\
 \frac{(\lambda x)Mx}{(\lambda xy)Mx} S \quad \frac{}{(\lambda xy)y} 0}{\frac{(\lambda xy)Mxy}{(\lambda x)\lambda y. Mxy} \lambda} @ \\
 \frac{(\lambda x)\lambda y. Mxy}{(\ ))M} \lambda \\
 \frac{(\ ))M}{(\ ))M} \text{FIX, } u
 \end{array}$$







# Proof system $\mathbf{Reg}^+$ for strong regularity

$$\frac{}{(\lambda \vec{x} y) y} 0 \quad \frac{(\lambda \vec{x} y) M_0}{(\lambda \vec{x}) \lambda y. M_0} \lambda$$

$$\frac{(\lambda \vec{x}) M_0 \quad (\lambda \vec{x}) M_1}{(\lambda \vec{x}) M_0 M_1} @$$

$$\frac{(\lambda x_1 \dots x_{n-1}) M}{(\lambda x_1 \dots x_n) M} \text{S (if the binding } \lambda x_n \text{ is vacuous)}$$

$$[(\lambda \vec{x}) M]^u$$

$$\frac{\mathcal{D}_0}{(\lambda \vec{x}) M} \text{FIX, } u \text{ (if } |\mathcal{D}_0| \geq 1)$$

# Proof system $\mathbf{Reg}_0^+$ for strong regularity

$$\frac{}{(\lambda \vec{x} y) y} 0 \qquad \frac{(\lambda \vec{x} y) M_0}{(\lambda \vec{x}) \lambda y. M_0} \lambda$$

$$\frac{(\lambda \vec{x}) M_0 \quad (\lambda \vec{x}) M_1}{(\lambda \vec{x}) M_0 M_1} @$$

$$\frac{(\lambda x_1 \dots x_{n-1}) M}{(\lambda x_1 \dots x_n) M} \text{S (if the binding } \lambda x_n \text{ is vacuous)}$$

$$[(\lambda \vec{x}) M]^u$$

 $\mathcal{D}_0$ 

$$\frac{(\lambda \vec{x}) M}{(\lambda \vec{x}) M} \text{FIX, } u \text{ (if } |\mathcal{D}_0| \geq 1, \text{ and, on threads in } \mathcal{D}_0 \text{ down from assumptions } ((\lambda \vec{x}) M)^u, \text{ prefix-lengths are } \geq |\vec{x}|)$$

# Proof system $\text{ann-Reg}_0^+$ for $\lambda_{\text{letrec}}$ -expressibility

$$\frac{}{(\lambda \vec{x} y) y : y} 0 \qquad \frac{(\lambda \vec{x} y) L : M}{(\lambda \vec{x}) \lambda y. L : \lambda y. M} \lambda$$

$$\frac{(\lambda \vec{x}) L_0 : M_0 \quad (\lambda \vec{x}) L_1 : M_1}{(\lambda \vec{x}) L_0 L_1 : M_0 M_1} @$$

$$\frac{(\lambda x_1 \dots x_{n-1}) L : M}{(\lambda x_1 \dots x_n) L : M} \text{S (if the binding } \lambda x_n \text{ is vacuous)}$$

$$[(\lambda \vec{x}) c_u : M]^u$$

$$\mathcal{D}_0$$

$$\frac{(\lambda \vec{x}) L[u := c_u] : M}{(\lambda \vec{x}) (\text{letrec } u = L \text{ in } u) : M} \text{FIX, } u$$

(if  $|\mathcal{D}_0| \geq 1$ , and, on threads in  $\mathcal{D}_0$  down from assumptions  $((\lambda \vec{x}) M)^u$ , **prefix-lengths are  $\geq |\vec{x}|$** )

# Proof system $\mathbf{Reg}^+$ for strong regularity

Example ( $M = \lambda xy. M x y$ )

$$\begin{array}{c}
 \frac{((\ )M)^u}{(\lambda x)M} \text{ S} \quad \frac{}{(\lambda x)x} \text{ 0} \\
 \hline
 (\lambda x)Mx \quad \text{0} \\
 \frac{}{(\lambda xy)Mx} \text{ S} \quad \frac{}{(\lambda xy)y} \text{ 0} \\
 \hline
 (\lambda xy)Mxy \quad \text{0} \\
 \frac{}{(\lambda x)\lambda y. Mxy} \lambda \\
 \frac{}{(\ )M} \lambda \\
 \frac{}{(\ )M} \text{ FIX, } u
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# Proof system $\mathbf{Reg}_0^+$ for strong regularity

Example ( $M = \lambda xy. Mxy$ )

$$\begin{array}{c}
 \frac{((\ )M)^u}{(\lambda x)M} \text{ S} \quad \frac{}{(\lambda x)x} 0 \\
 \hline
 (\lambda x)Mx \quad \textcircled{\text{C}} \\
 \frac{(\lambda x)Mx}{(\lambda xy)Mx} \text{ S} \quad \frac{}{(\lambda xy)y} 0 \\
 \hline
 (\lambda xy)Mxy \quad \textcircled{\text{C}} \\
 \frac{(\lambda xy)Mxy}{(\lambda x)\lambda y. Mxy} \lambda \\
 \frac{(\lambda x)\lambda y. Mxy}{(\ )M} \lambda \\
 \frac{(\ )M}{(\ )M} \text{ FIX, } u
 \end{array}$$

# Proof system $\text{ann-Reg}_0^+$ for $\lambda_{\text{letrec}}$ -expressibility

## Example

$$\frac{\frac{((\ ) c_u : M)^u}{(\lambda x) c_u : M} \text{S} \quad \frac{}{(\lambda x) x : x} 0}{\frac{(\lambda x) c_u x : Mx}{(\lambda xy) c_u x : Mx} \text{S} \quad \frac{}{(\lambda xy) y : y} 0}{} @}{\frac{(\lambda xy) c_u xy : Mxy}{(\lambda x) \lambda y. c_u xy : Mxy} \lambda}{\frac{(\ ) \lambda xy. c_u xy : \lambda xy. Mxy}{(\ ) (\text{letrec } u = \lambda xy. uxy \text{ in } u) : M} \text{FIX, } u} \lambda$$

# Main Result

## Theorem

*An infinite  $\lambda$ -term is  $\lambda_{\text{letrec}}$ -expressible if and only if it is strongly regular.*

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## Proof.

$\Leftarrow$ : For every infinite  $\lambda$ -term  $M$ :

$$\begin{aligned}
 M \text{ is strongly regular} &\implies \vdash_{\mathbf{Reg}^+} () M \\
 &\implies \vdash_{\mathbf{Reg}_0^+} () M \\
 &\implies \vdash_{\mathbf{ann-Reg}_0^+} () L : M \\
 &\implies L \rightarrow_{\nabla}^{\omega} M \\
 &\implies L \text{ expresses } M
 \end{aligned}$$



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 \end{aligned}$$

- $\Rightarrow$ :
- ▶ define:  $\mathbf{Reg}^+$ -observations (& observ. graphs) for  $\lambda_{\text{letrec}}$ -terms
  - ▶ show: for every  $\lambda_{\text{letrec}}$ -term  $L$ , the  $\mathbb{S}_{\text{eag}}^+$ -observ. graph of  $L$  is finite
  - ▶ show: observation graph of  $\lambda_{\text{letrec}}$ -term  $L$  is finite
    - $\Rightarrow$  observation graph of the unfolding  $\mathcal{U}(L)$  of  $L$  is finite.

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# Main theorem (extended)

## Theorem

For all infinite  $\lambda$ -terms the following are equivalent:

- (i)  $M$  is  $\lambda_{\text{letrec}}$ -expressible.
- (ii)  $M$  is strongly regular.
- (iii)  $M$  only contains *finite binding-capturing chains*.