

**INTRODUCTION TO MODAL LOGIC 2016
HOMEWORK 5**

- Deadline: November 29 — at the **beginning** of class.
- Grading is from 0 to 100 points.
- Results from the exercise class may be used in the proofs
- Success!

(1) (30pt)

- (a) Prove that $\mathbf{S5} = \mathbf{K4} + (\Box p \rightarrow p, p \rightarrow \Box \Diamond p)$ is sound and complete with respect to the class of frames (W, R) , where R is an equivalence relation (i.e., a reflexive, transitive and symmetric relation).
- (b) Use (a) to show that $\mathbf{S5}$ is sound and complete with respect to the class of frames (W, R) where for each $w, v \in W$ we have Rwv . (Hint: use point-generated submodels, Prop 2.6 from the book.)

(2) (30pt) (From the 2014 Exam) In the following exercise you can use that the canonical model for $\mathbf{S4.3}$ is reflexive and transitive.

- (a) Show that the canonical model for the modal logic

$$\mathbf{S4.3} = \mathbf{S4} + \Box(\Box p \rightarrow q) \vee \Box(\Box q \rightarrow p)$$

has no branching to the right. Recall that a reflexive Kripke frame has no branching to the right if

$$\forall x \forall y \forall z ((Rxy \wedge Rxz) \rightarrow (Ryz \vee Rzy)).$$

You are not allowed to use Sahlqvist completeness theorem.

- (b) Deduce that $\mathbf{S4.3}$ is sound and complete with respect to reflexive transitive frames with no branching to the right.

(3) (20pt) (Item (a) is from the 2014 Exam)

- (a) Prove that for any modal formulas φ and ψ we have

$$\vdash_{\mathbf{K}} \Box \varphi \vee \Box \psi \text{ implies } \vdash_{\mathbf{K}} \varphi \text{ or } \vdash_{\mathbf{K}} \psi.$$

(Hint: use completeness of \mathbf{K} with respect to Kripke frames.)

- (b) Show that the above property does not hold for all normal modal logics. That is, give an example of a normal modal logic L which does not satisfy it.

- (4) (20pt) Given a frame class \mathcal{C} , let $\Theta(\mathcal{C}) = \text{Log}(\mathcal{C})$ and given normal modal logic L let $\text{Fr}(L)$ be the class of frames where L is valid.
- (a) What does it mean for a logic L if $L = \Theta(\text{Fr}(L))$? Give an example of a logic (modal or temporal) for which it does not hold.
- (b) What does it mean for a frame class \mathcal{C} if $\mathcal{C} = \text{Fr}(\Theta(\mathcal{C}))$? Give an example of a frame class \mathcal{C} for which it does not hold.